



Drilling Into Drill: Flexiperf

Jacques Nadeau, Architect and VP Apache Drill February 19, 2015



© 2014 MapR Technologies



"Drill isn't just about SQL-on-Hadoop. It's about SQL-onpretty-much-anything, immediately, and without formality." -Andrew Brust, GigaOM Research, Dec 2014





Agenda

- What?
 - SQL like mom made
 - Punk SQL
- How?
 - Flexibility
 - Performance
- Who & When



SQL Like Mom Made

SoH Table Stakes: Warehousing and Business Intelligence

ANSI Syntax

- SELECT, FROM, WHERE, JOIN, HAVING, ORDER BY, WITH, CTAS, OVER*, ROLLUP*, CUBE*, ALL, EXISTS, ANY, IN, SOME
- VarChar, Int, BigInt, Decimal, VarBinary, Timestamp, Float, Double, etc.
- Subqueries, scalar subqueries*, partition pruning, CTE

Interactive SQL Workloads

- Data warehouse offload
- Tableau, ODBC, JDBC
- TPC-H & TPC-DS-like workloads

Standard Hadoop Tools

- Supports Hive SerDes
- Supports Hive UDFs
- Supports Hive Metastore

Punk SQL

EA.

Punk SQL: SQL for a Hadoop World New Workloads

- Path based queries and wildcards
 - select * from /my/logs/
 - select * from /revenue/*/q2
- Modern data types
 - Any, Map, Array (JSON)
- Complex Functions and Relational Operators
 - FLATTEN, kvgen, convert_from, convert_to, repeated_count, etc

- JSON Sensor analytics
- Complex data analysis
- Alternative DSLs

New Ways to Work

- Query without prep
- Workspaces without admin intervention
- Expose query as MapReduce
- Expose query as Spark RDD





How?





Flexibility your tool should be flexible...

adia

so you don't have to be

London 2012

8811389

saypo

Flexibility is a Vision, Usability, a Religion

- Deployment
- Data Model
- Schema
- Security
- Access Methods



Supporting the Changing Roles of Big Data

Data Dev Circa 2000

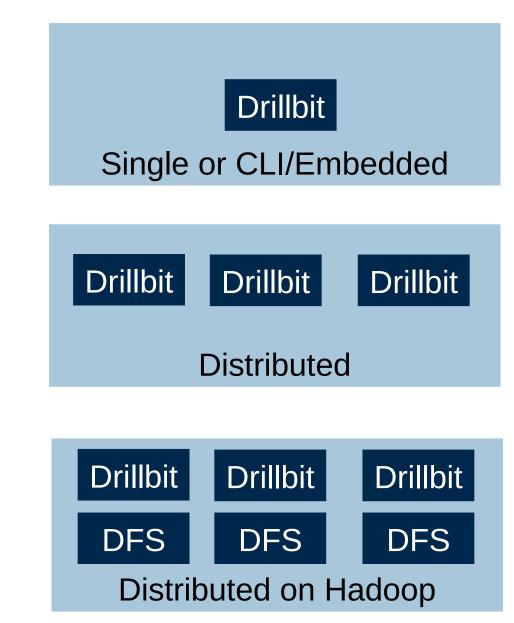
- 1. Developer comes up with requirements
- 2. DBA defines tables
- 3. DBA defines indices
- 4. DBA defines FK relationships
- 5. Developer stores data
- 6. BI builds reports
- 7. Analyst views reports
- 8. DBA adds materialized views

Data Today

- Developer builds app, defines schema, stores data
- 2. Analyst queries data
- 3. Data engineer fixes performance problems or fills functionality gaps

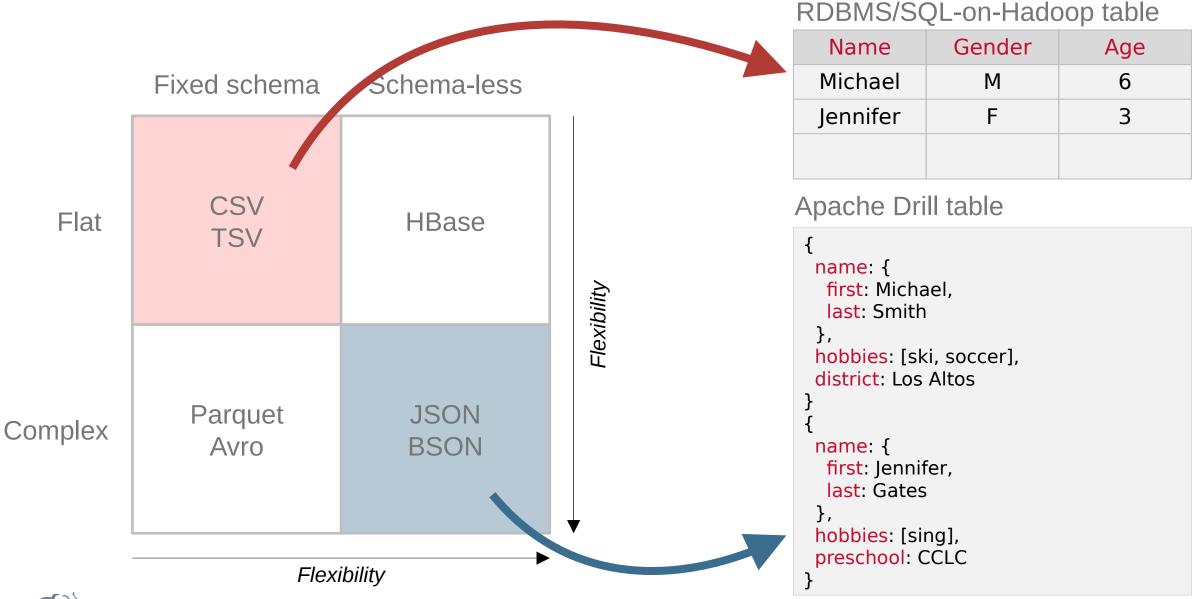
Drill Deployment is Easy

- Single Daemon for all purposes
- No special considerations for scaling or availability
- With or without DFS
- Works with other data systems (Mongo, Cassanrda & JDBC coming soon)
- Runs on Linux, Mac or Windows
- No separate database
- JSON everything
- Access via HTTP, Java, C, JDBC, ODBC, CLI





Drill Provides A Flexible Data Model



13

Leave Your Data Where it is. Access it Centrally & Uniformly.

- Drill is storage agnostic
- Interacts to storage through
 plugins
- Storage plugins expose optimizer rules
 - Optimizer rules work directly on logical operation to expose maximum capabilities
- Reference multiple Hive, HBase, MongoDB, DFS, etc systems

Apache Drill		
Enabled Storage Plugi	ins	
ср	Update	Disable
dfs	Update	Disable
fs	Update	Disable
hbase-prod	Update	Disable
hive-marketing	Update	Disable
mongo-promotions	Update	Disable



Leverage that Massive Scalable Redundant Infrastructure

Single Store for Data and Metadata

- HDFS is already your single canonical store
- Don't create a secondary metadata store

Avoid Metadata Management and synchronization

- Store metadata inline
 - If you can't, store it next to files
- Move directories around at will
- Delete things at will

Flexibility in how you describe your data

- Drill doesn't require schema, detects file types based on
 - extensions
 - magic bytes (e.g. PAR1)
 - systems settings
- Query can be planned on any file, anywhere
- Data types are determined as data arrives
- Some formats have known schema
 - If they don't, you can expose them as such through views
 - Views are simply JSON files that define view SQL

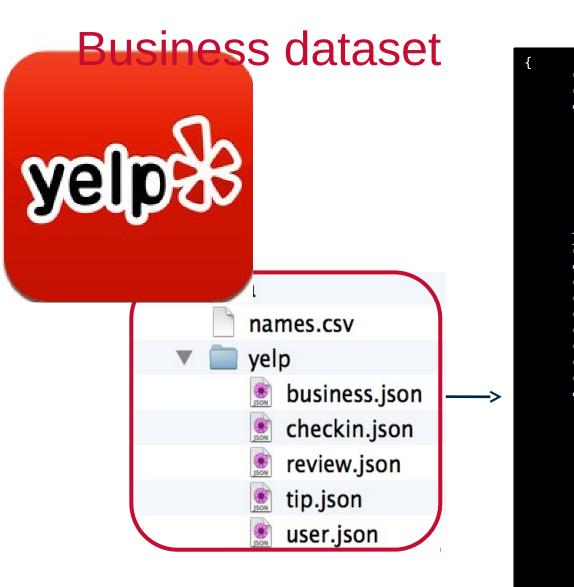




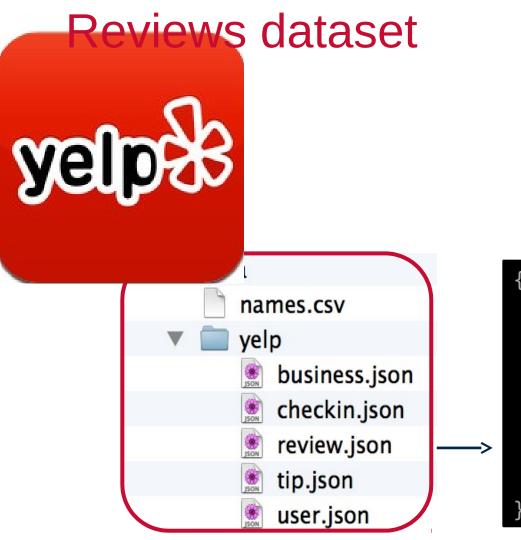
Product Walkthrough







"business id": "4bEj0 yTaDG24SY5TxsaUNQ", "full address": "3655 Las Vegas Blvd S\nThe Strip\nLas Vegas, NV 89109", "hours": { "M onday": { "close": "23:00", "open": "07:00"}, "Tuesday": { "close": "23:00", "open": "07:00"}, "Friday": { "close": "00:00", "open": "07:00"}, "W ednesday": { "close": "23:00", "open": "07:00"}, "Thursday": { "close": "23:00", "open": "07:00"}, "Sunday": { "close": "23:00", "open": "07:00"} , "Saturday": { "close": "00:00", "open": "07:00"} }, "open": true. "categories": ["Breakfast & Brunch", "Steakhouses", "French", "Restaurants"], "city": "Las Vegas", "review count":4084, "nam e": "M on Am iGabi", "neighborhoods": ["The Strip"], "longitude": -115.172588519464, "state":"NV", "stars": 4.0, "attributes": { "Alcohol": "full bar", "Noise Level": "average", "Has TV": false, "Attire": "casual", "Am bience":{ "rom antic": true, "intim ate": false, "touristy": false, "hipster":false, "classy": true, "trendy": false, "casual": false }, "Good For": { "dessert": false, "latenight": false, "lunch": false, "dinner": true, "breakfast": false, "brunch": false},



```
"votes": { "funny": 0, "useful": 2, "cool": 1 },
"user_id": "Xqd0D zHaiyRqVH3W RG7hzg",
"review_id": "15SdjuK7Dm YqUAj6rj6owg",
"stars": 5,
"date": "2007-05-17",
"text": "dr.goldberg of @rs everything ...",
"type": "review",
"business_id": "vcNAW iLM 4dR7D 2nww J7nCA"
```

Zero to Results in 2 minutes

<pre>\$ tar -xvzfapache -d rill-0 .7 .0 .tar.gz</pre>	Install	
\$bin/sqlline-ujdbc:drill:zk=local	Launch shell (em bedded m ode)	
> SELECT state, city, count(*) AS businesses		
FROM dfs.yelp.`business.json`		
GROUP BY state, city ORDER BY businesses DESC LIM IT 10;	Query files and directories	
++++ state city businesses +++		
NV Las Vegas 12021 AZ Phoenix 7499 AZ Scottsdale 3605 EDH Edinburgh 2804		
AZ Mesa 2041 AZ Tempe 2025	- Results	
NV Henderson 1914 AZ Chandler 1637	- Resulus	
WI Madison 1630 AZ Gendale 1196		

Intuitive SQL access to complex data

//It's Friday 10pm in Vegas and Looking for Hummus

> SELECT nam e, stars, b.hours.Friday friday, categories FROM dfs.yelp.`business.json`b W HERE b.hours.Friday.`open` < '22:00'AND b.hours.Friday.`close` > '22:00'AND <u>REPEATED_CONTAINS</u> (categories, M editerranean') AND city = 'Las Vegas' ORDER BY stars DESC LIMIT 2;

Query data with any levels of nesting

| nam e | stars | friday | categories | + -----+ ----+ -----+ -----+ |0 lives |4.0 |{"close":"22:30","open":"11:00"} |["M editerranean","Restaurants"] | |M arrakech M oroccan Restaurant |4.0 |{"close":"23:00","open":"17:30"} | ["M editerranean","M iddle Eastern","M oroccan","Restaurants"] | + -----+ -----+ -----+

ANSI SQL compatibility

//Get top cool rated businesses

SELECT b.nam e from dfs.yelp.`business.json` b W HERE b.business_id IN (SELECT r.business_id FROM dfs.yelp.`review.json` r GROUP BY r.business_id HAVING SUM (r.votes.cool) > 2000 ORDER BY SUM (r.votes.cool) DESC);

Use fam iliar SQL functionality (Joins, Aggregations, Sorting, Subqueries, SQL data types)

	_

| nam e

+ -----4

|EarlofSandwich | |XS Nightclub | |The Cosm opolitan ofLas Vegas | |Wicked Spoon | + ------+

Logical views

//Create a view combining business and reviews datasets

> CREATE OR REPLACE VIEW dfs.tm p.BusinessReviews AS SELECT b.nam e, b.stars, r.votes.funny, r.votes.useful, r.votes.cool, r.`date` FROM dfs.yelp.`business.json` b, dfs.yelp.`review.json` r W HERE r.business id = b.business id;

> SELECT COUNT(*) AS TotalFROM dfs.tm p.BusinessReviews;

+----+ | Total | +----+ | 1125458 | +----+ Lightweightfle system based views for granular and decentralized data m anagem ent

Materialized Views AKA Tables

> ALTER SESSION SET `store.form at` = 'parquet';

Save analysis results as tables using fam iliar CTAS syntax

+	-++
Fragm e	nt Num ber of records w ritten
+	-++
1_0	176448
1_1	192439
1_2	198625
1_3	200863
1_4	181420
1_5	175663
+	++



Working with repeated values





Extensions to ANSI SQL to work with repeated values

//Flatten repeated categories

> SELECT nam e, categories FROM dfs.yelp.`business.json` LIM IT 3;

+ -----+ ----+ | name |categonies| + -----+ |Enic Goldberg, MD | ["Doctors", "Health & Medical"] | |Pine Cone Restaurant | ["Restaurants"] | |Deforest Family Restaurant | ["Amenican (Traditional)", "Restaurants"] | + ------+

> SELECT nam e, FLATTEN (categories) AS categories FROM dfs.yelp.`business.json` LIM IT 5;

name | categories

|Eric Goldberg, MD |Doctors | |Eric Goldberg, MD |Health & Medical| |Pine Cone Restaurant |Restaurants | |Deforest Fam ily Restaurant |American (Traditional) | |Deforest Fam ily Restaurant |Restaurants |

Dynam ically fatten repeated and nested data elem ents as part of SQ L queries. No ETL necessary

Extensions to ANSI SQL to work with repeated values

//Getm ost com m on business categories

> SELECT category, count(*) AS categorycount
FROM (SELECT nam e, FLATTEN (categories) AS category
FROM dfs.yelp.`business.json`) c
GROUP BY category ORDER BY categorycountDESC;

```
+ -----+

| category |categorycount|

+ -----+

|Restaurants |14303 |

...

|Australian |1 |

|Boat Dealers |1 |

|Firew ood |1 |

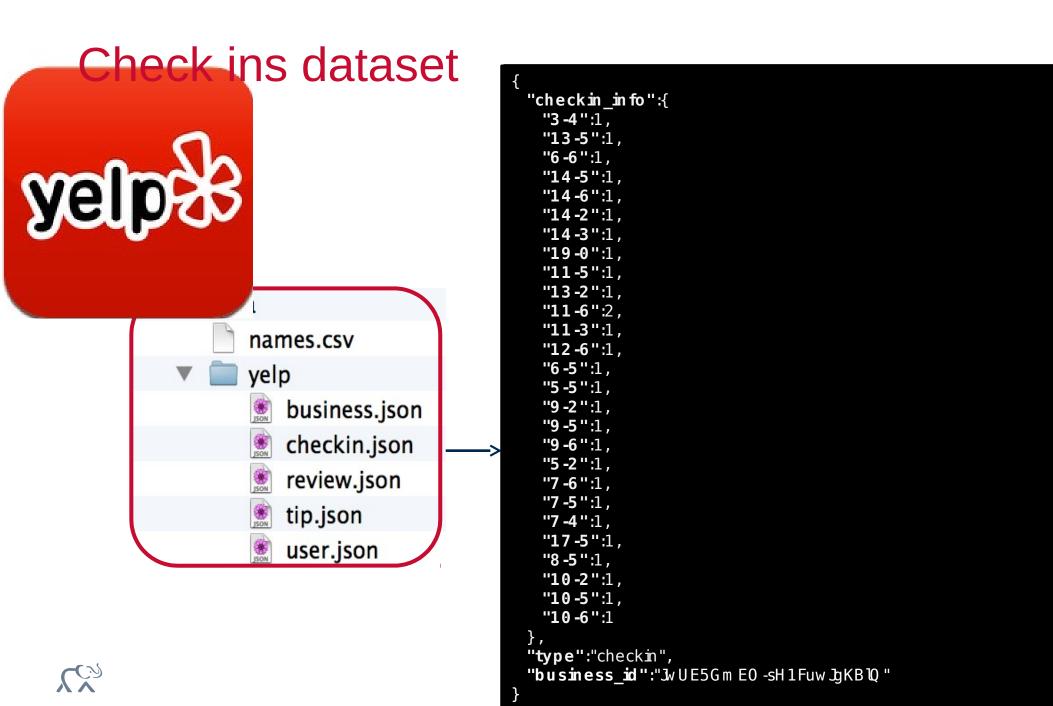
+ -----+
```



Working with Dynamic Columns







Makes it easy to work with dynamic/unknown columns

> jdbc:drill:zk= local> SELECT KVGEN (checkin_info) checkins FROM dfs.yelp.`checkin.json` LIM IT 1;

+ -----+

| checkins

+ -----+

> jdbc:drill:zk= local> SELECT FLATTEN (KVGEN (checkin_info)) checkins FROM dfs.yelp.`checkin.json`limit6;

-----+ checkins { "key":"3-4","value":1} | |{"key":"13-5","value":1} | |{"key":"6-6","value":1} { "key":"14-5","value":1} { "key":"14-6","value":1} { "key":"14-2","value":1 }

⊬ ----≁

Convert Map with a wide set of dynam ic colum ns into an array of key-value pairs

Makes it easy to work with dynamic/unknown columns

//Count totalnum berof check ins on Sunday m idn ight

 jdbc:drill:zk= local> SELECT SUM (checkintblcheckins.`value`) as SundayM idnightCheckins FROM
 (SELECT FLATTEN (KVGEN (checkin_info)) checkins FROM dfs.yelp.checkin.json`) checkintbl
 W HERE checkintbl.checkins.key= '23-0';

+ -----+

|SundayMidnightCheckins|

+ -----+

|8575

+ -----+

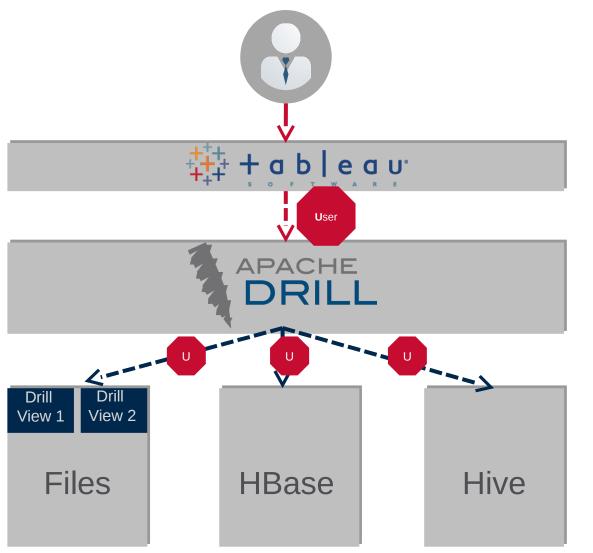


Secure Access





Access control that scales

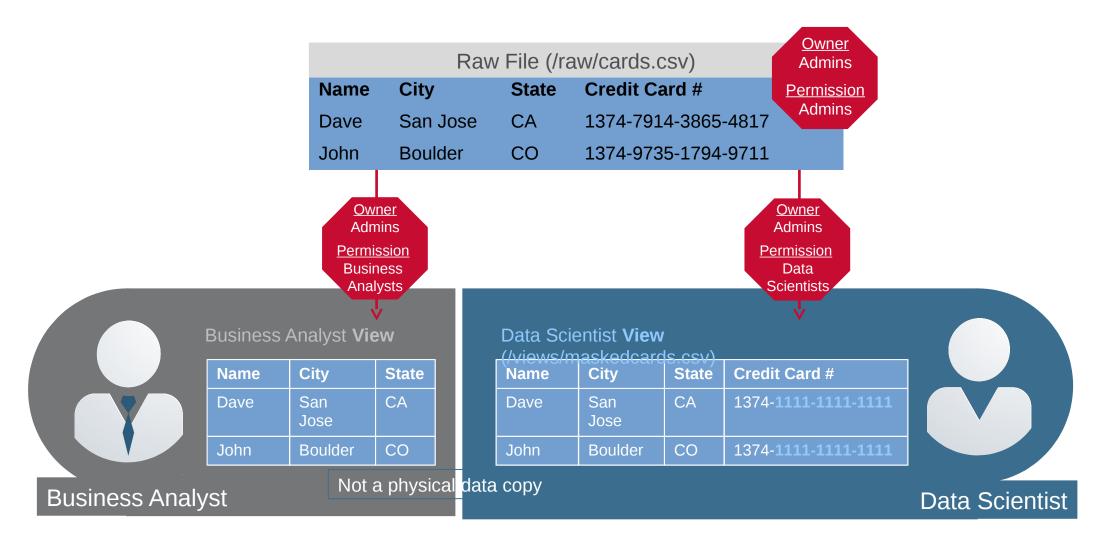




Fine-grained row and column level access control with Drill Views – no centralized security repository required

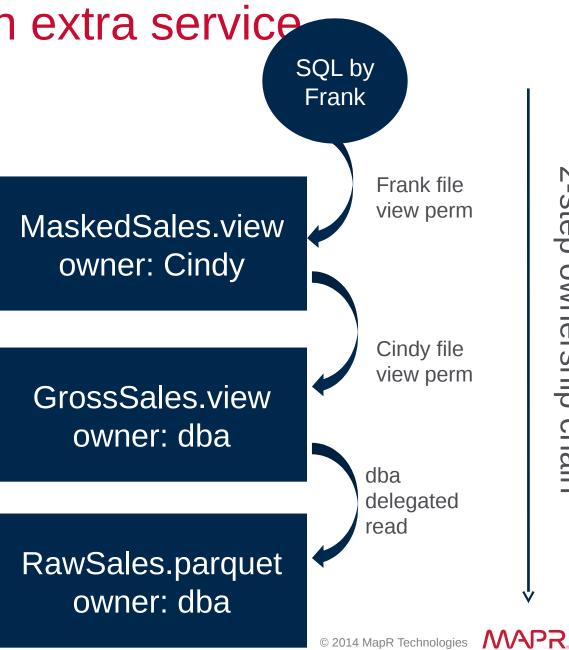
 $\mathbf{V}_{\mathcal{O}}$

Granular security permissions through Drill views



Secure your data without an extra service

- Drill Views
- Ownership chaining with configurable delegation TTL
- Leverages existing HDFS ACLs
- Complete security solution without additional services or software



لاريم

2-step ownership chair

Summary

- Logical
 - No physical data copies/silos
- Granular
 - Row level and column level security controls
- De-centralized
 - User impersonation respecting storage system permissions
 - No separate permission repository for granular controls
 - Integrated with Hadoop File System permissions and LDAP
- Self-service w/ governance
 - If you have access to data, you control who and how widely can access it
 - Audits

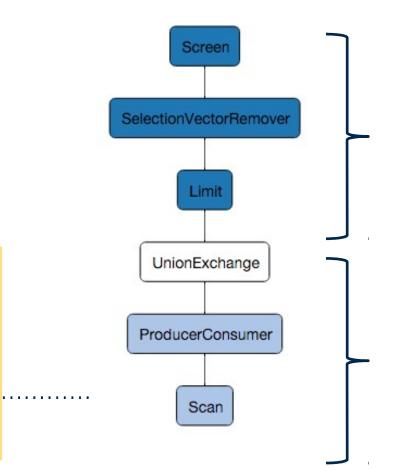
Core Drill Architectural Goals

- Go fast when you don't know anything
 - And do "the right thing"
- Go faster when you do know things

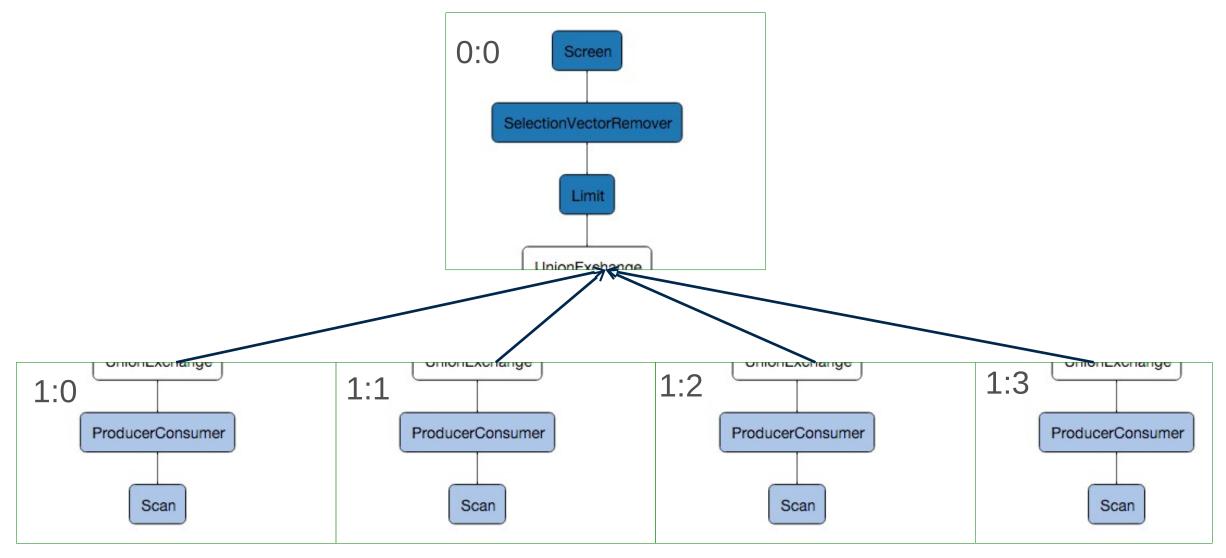
An Optimistic, Pipelined Purpose-Built DAG Engine

- Three-Level DAG
- Major Fragments (phases)
- Minor Fragments (threads)
- Operators (in-thread operations)

> explain plan for select * from customer limit 5;	
00-00	Screen
00-01	SelectionVectorRemover
00-02	Limit(fetch=[5])
 .00-03	UnionExchange
01-01	ProducerConsumer
01-02	Scan(groupscan=[ParquetGroupScan [



Each phase (MajorFragment) gets Parallelized (MinorFragment)



Reading Data Quickly, Moving Data Quickly

- Highly Optimized Native Drill Readers:
 - Vectorized Parquet, Text/CSV, JSON
 - Also works with all Hive supported formats
- Drill supports partition pruning
 - Adding direct physical property exposure soon for highly optimized cases
- Drill parallelizes to maximum level format allows
 - Also balances data locality and maximum parallelization
- Bespoke Asynchronous Zero-Copy RPC Layer
 - Built specifically for Drill's internal data format

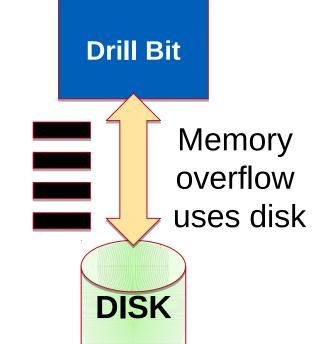
Parquet: The Format for the next decade

Apache Drill & Apache Parquet communities working together

- Better Ecosystem integration and decentralized metadata
 - Self Description capabilities
 - Ecosystem support for logical data types
- Enhanced Performance
 - New vectorized reading
 - Enhanced memory pooling and management
 - Indexing*
 - Better metadata positioning (for improved page pruning)*
 - Enhanced vectorization and late materialization reading*

Value Vectors & Record Batches: Drill's In-memory Columnar Work Units

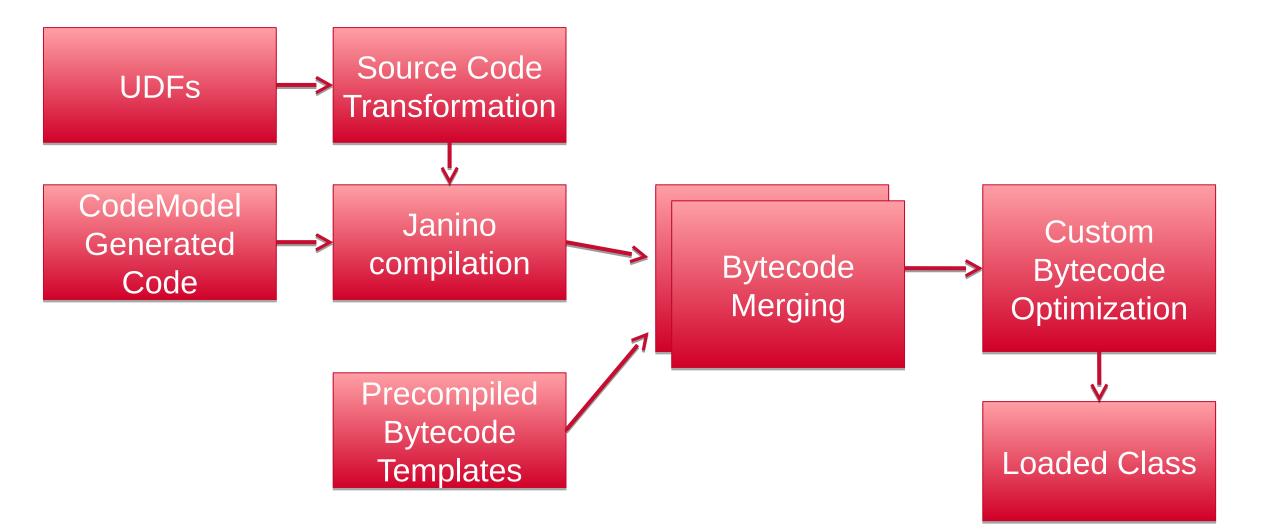
- Random access: sort without copy or restructuring
- Fully specified in memory shredded complex data structures
- Remove serialization or copy overhead at node boundaries
- Spool to disk as necessary
- Interact with data in Java or C without copy or overhead



Runtime Compilation and Multi-phased planning

- Drill does best effort initial query parsing, planning and validation
 - Where Drill doesn't understand data, it provides support for ANY type, allowing late type binding.
- At execution time, individual nodes do secondary pass
 - Schema <--> Query parsing and validation
 - Type casting, coercion and promotion
 - Compilation based on schema requirements
- As schema changes, Drill supports recompilation of each operator as necessary

Runtime Compilation Pattern



Advanced Compilation Techniques

- Optimization based on observation and assembly
- Drill does a number of pre-machine-code-compilation optimizations to ensure efficient execution
- Some examples:
 - Removal of type and bounds checking
 - Direct micro pointers for in-record-batch references
 - Little endian data formats
 - Bytecode-level scalar replacement



Drill Does Vectorization & Supports Columnar Functions

- Drill often operates on more than one record at a time
 - Word-sized manipulations
 - SIMD instructions
 - Manually coded algorithms
- Columnar Functions Improve Many Operations
 - Bitmaps allow lightning fast null-checks and reduction in branching
 - Type demotion to reduce memory and CPU computation overhead
 - Direct Conversions where possible

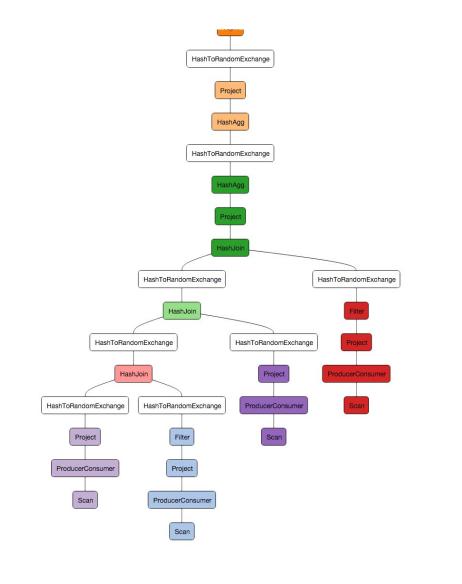


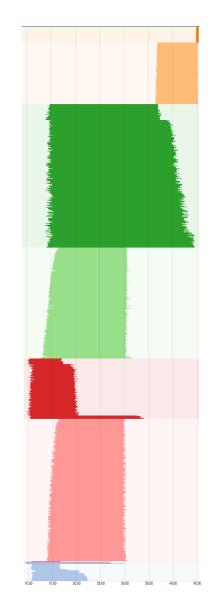
Drill provides advanced query telemetry

- What happened during query for all three levels of DAG execution
- Each profile is stored as JSON file for easy review, sharing and backup (at end of query execution)
- Profiles can be analyzed using Drill, allows:
 - easy longitudinal analysis of workload
 - multi-tenancy performance analysis
 - impact of configuration changes to benchmark workloads



A Color-coded visual layout and Gant timing chart is provided





Memory Efficiency

- Drill's in memory representation is designed to minimize memory overhead
- Custom implementation of columnar-aware data structures including hash tables, sort operations, etc.
 - For example, entry overhead for hash table is 8 bytes per value
 - Sort pointer overhead for sort is 2 to 4 bytes per entry
- Adding support for compressed columnar representation further to improve compactness

Scale and Concurrency

- Drill's execution model leverages both local and remote exchanges for changes in parallelization
 - Muxing Local, Demuxing Local, Broadcast, Hash to Merge, Hash to Random, Ordered Partition, Single Merge, Union
- All operations can be parallelized at the thread and node level
- Thread count and parallelization are influenced by data size, query phasing and system load
 - Administrators have basic queuing control to manage workload
- All threads are independent and pipelined, all run in a single process per node
- Node <> node communication is multiplexed, push-based with sub-socket backpressure support
- Testing has proceeded up to 150 nodes. Target is 1000 nodes by GA.
 - Drill adapts data transfer size, buffers, muxing and other operations based on query scale to minimize n² multiplier effects